

# BBC Communications in GNURadio



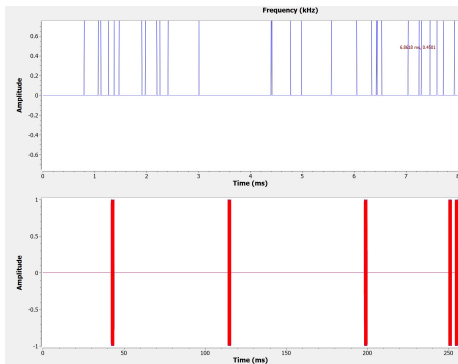
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# OVERVIEW

- BBC Overview
- GNURadio implementation
- Lessons learned
- Future Development



# WHAT IS BBC?

- Keyless
- Concurrent-codes
- Spread spectrum
- Jam-resistant (not LPI)
- Invertable hash algorithm encoding
- Limitation: Low throughput
- Limitation: modulation type
- Limit: Mark Density

No pre-shared key

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Can encode multiple messages  
in a single codeword

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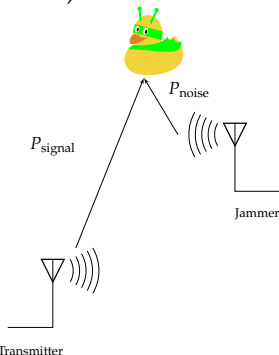
Transmitted signal size greater than message

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→ Cannot remove information that exists

→ Additional information in codeword can be ignored (to an extent)



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# WHAT IS BBC?

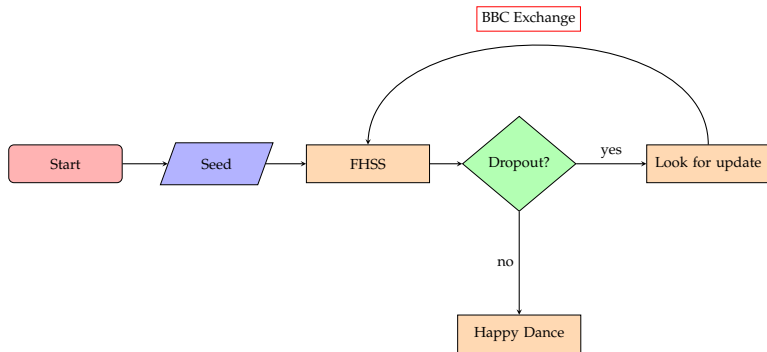
- Keyless
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- **Limit: Mark Density**

## Mark Density

$$\mu = \frac{n_{\text{bits, message}}}{n_{\text{bits, codeword}}}$$

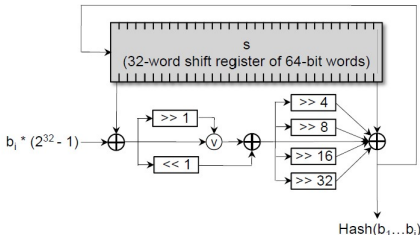
Limit  $\rightarrow 0.5$

# USE CASE - FHSS SEQUENCE



# ENCODING – THE GLOWWORM HASH

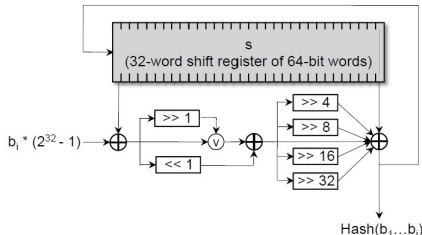
- Input is length of codeword
- Output is “sparse” codeword
- Shift register produces mark locations in codeword
- Shift register values depend on current substring
- Example →  
Message = 0011 0010  
Codeword = 16bits



<sup>1</sup> Baird, Carlisle, Bahn. The Glowworm Hash: Increased Speed and Security for BBC Unkeyed Jam Resistance. 2012.

# ENCODING – THE GLOWWORM HASH

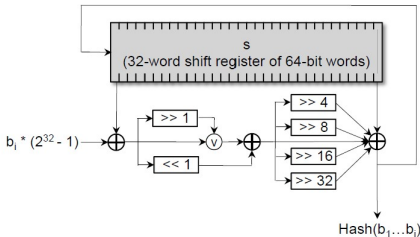
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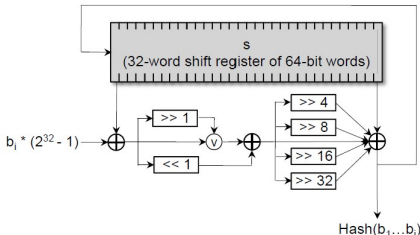
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- Example →  
 Message = 0011 0010  
 Codeword = 16bits

## MESSAGE:

0	0	1	1	0	0	0	0
0	1	2	3	4	5	6	7

## GLOWWORM HASH:

## CODEWORD:

0	0	0	0	0	0	0	0
0	1	2	3	4	5	6	7...
0	0	0	0	0	0	0	0
8	9	1	11	12	13	14	15

# ENCODING – THE GLOWWORM HASH



- Input is length of codeword
- Output is “sparse” codeword
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- Shift register values depend on current substring
- Example →  
 Message = 0011 0010  
 Codeword = 16bits

**MESSAGE:**

0	0	1	1	0	0	0	0
0	1	2	3	4	5	6	7

**GLOWWORM HASH:** 5

**CODEWORD:**

0	0	0	0	1	0	0	0
0	1	2	3	4	5	6	7...
0	0	0	0	0	0	0	0
8	9	10	11	12	13	14	15



# ENCODING – THE GLOWWORM HASH



- Input is length of codeword
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- Example →  
 Message = 0011 0010  
 Codeword = 16bits

**MESSAGE:**

0	0	1	1	0	0	0	0
0	1	2	3	4	5	6	7

**GLOWWORM HASH:** 0

**CODEWORD:**

1	0	0	0	1	0	0	0
0	1	2	3	4	5	6	7...
0	0	0	0	0	0	0	0
8	9	10	11	12	13	14	15

# ENCODING – THE GLOWWORM HASH



- Input is length of codeword
- Output is “sparse” codeword
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- Shift register values depend on current substring
- Example →  
 Message = 0011 0010  
 Codeword = 16bits

**MESSAGE:**

0	0	1	1	0	0	0	0
0	1	2	3	4	5	6	7

**GLOWWORM HASH:** 14

**CODEWORD:**

1	0	0	0	1	0	0	0
0	1	2	3	4	5	6	7...
0	0	0	0	0	0	1	0
8	9	1	11	12	13	14	15

# ENCODING – THE GLOWWORM HASH



- Input is length of codeword
- Output is “sparse” codeword
- Shift register produces mark locations in codeword
- Shift register values depend on current substring
- Example →  
 Message = 0011 0010  
 Codeword = 16bits

**MESSAGE:**

0	0	1	1	0	0	0	0
0	1	2	3	4	5	6	7

**GLOWWORM HASH: 2**

**CODEWORD:**

1	0	1	0	1	0	0	0
0	1	2	3	4	5	6	7...
0	0	0	0	0	0	1	0
8	9	10	11	12	13	14	15

# ENCODING – THE GLOWWORM HASH



- Input is length of codeword
- Output is “sparse” codeword
- Shift register produces mark locations in codeword
- Shift register values depend on current substring
- Example →  
 Message = 0011 0010  
 Codeword = 16bits

**MESSAGE:**

0	0	1	1	0	0	0	0
0	1	2	3	4	5	6	7

**GLOWWORM HASH:** 0

**CODEWORD:**

1	0	1	0	1	0	0	0
0	1	2	3	4	5	6	7...
0	0	0	0	0	0	1	0
8	9	10	11	12	13	14	15

# ENCODING – THE GLOWWORM HASH



- Input is length of codeword
- Output is “sparse” codeword
- Shift register produces mark locations in codeword
- Shift register values depend on current substring
- Example →  
 Message = 0011 0010  
 Codeword = 16bits

**MESSAGE:**

0	0	1	1	0	0	0	0
0	1	2	3	4	5	6	7

**GLOWWORM HASH:** 0

**CODEWORD:**

1	0	1	0	1	0	0	0
0	1	2	3	4	5	6	7...
0	0	0	0	0	0	1	0
8	9	10	11	12	13	14	15

# ENCODING – THE GLOWWORM HASH



- Input is length of codeword
- Output is “sparse” codeword
- Shift register produces mark locations in codeword
- Shift register values depend on current substring
- Example →  
 Message = 0011 0010  
 Codeword = 16bits

**MESSAGE:**

0	0	1	1	0	0	0	0
0	1	2	3	4	5	6	7

**GLOWWORM HASH:** 3

**CODEWORD:**

1	0	1	1	1	0	0	0
0	1	2	3	4	5	6	7...
0	0	0	0	0	0	1	0
8	9	1	11	12	13	14	15

# ENCODING – THE GLOWWORM HASH



- Input is length of codeword
- Output is “sparse” codeword
- Shift register produces mark locations in codeword
- Shift register values depend on current substring
- Example →  
 Message = 0011 0010  
 Codeword = 16bits

**MESSAGE:**

0	0	1	1	0	0	0	0
0	1	2	3	4	5	6	7

**GLOWWORM HASH:** 12

**CODEWORD:**

1	0	1	1	1	0	0	0
0	1	2	3	4	5	6	7...
0	0	0	0	1	0	1	0
8	9	10	11	12	13	14	15

# ENCODING – THE GLOWWORM HASH

- Input is length of codeword
- Output is “sparse” codeword
- Shift register produces mark locations in codeword
- Shift register values depend on current substring

- Example →

Message = 0011 0010

Codeword = 16bits

- $\mu = \frac{b_{\text{message}}}{b_{\text{codeword}}} = \frac{6}{16} = 0.375$

**MESSAGE:**

0	0	1	1	0	0	0	0
0	1	2	3	4	5	6	7

**GLOWWORM HASH:**

**CODEWORD:**

1	0	1	1	1	0	0	0
0	1	2	3	4	5	6	7...
0	0	0	0	1	0	1	0
8	9	1	11	12	13	14	15



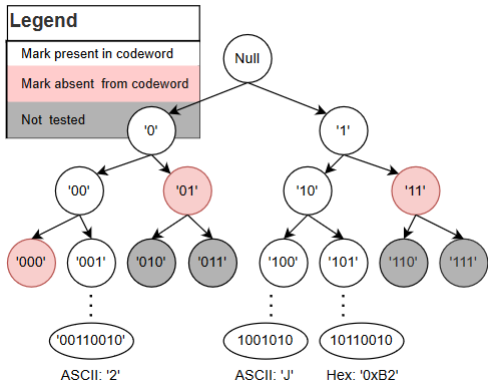
# DECODING

## • Depth-First Search:

- Inuitive to implement,
- Stack depth
- Inefficient use of memory

## • Breadth-First Search

- More difficult to implement
- Memory / CPU efficient



# JAM-RESISTANCE

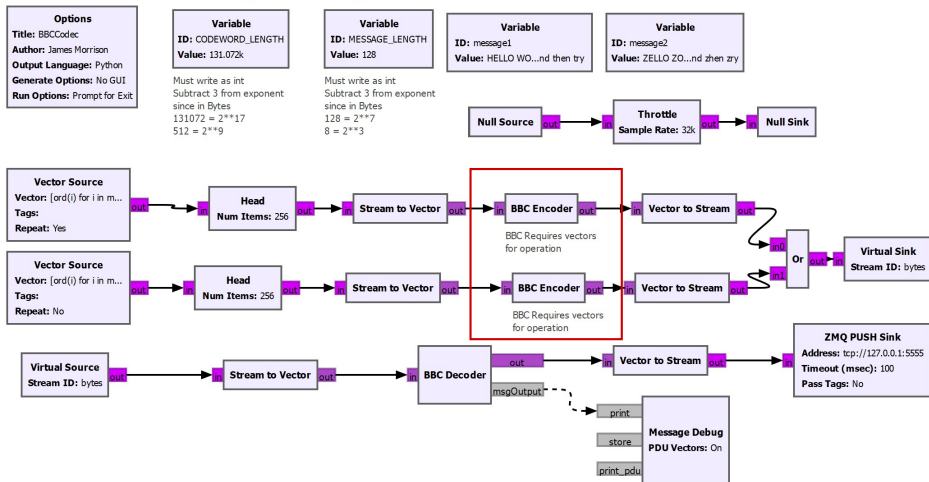


Jamming acts as a bit-wise “OR”

Origin	ASCII Character	BBC Codeword
Original message (TX)	2	10011100 00001010
Jamming signal	!	01000001 10001101
<b>Channel packet (OR)</b>	<b>N/A</b>	<b>11011101 10001111</b>

- Broadband noise attack → enough energy wins (always)
- Random marks add additional messages, don't affect existing
- Checksum bits protect against hallucinations

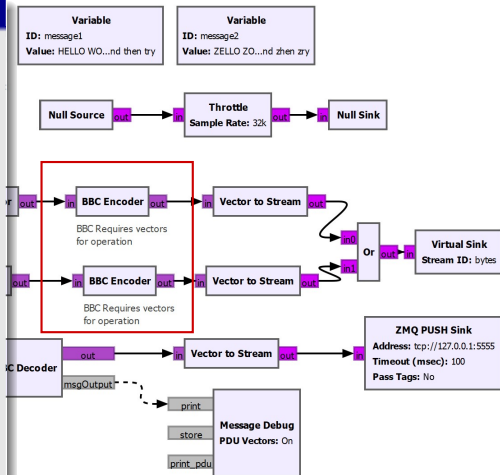
# BBC IN GNURADIO - ENCODE/DECODE



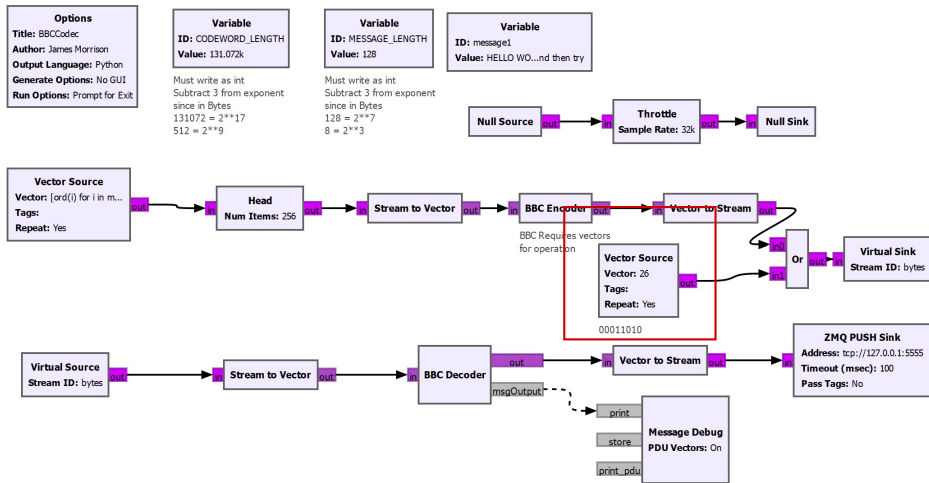
# BBC IN GNURADIO - ENCODE/DECODE

## Details

- Encoder
  - Input = bytes vectors of length MESSAGE LENGTH
  - Output = bytes vectors of length CODEWORD LENGTH
- Decoder
  - Input = bytes vectors of length CODEWORD LENGTH
  - Output = bytes vectors of length MESSAGE LENGTH



# BASIC TESTING – PACKET DENSITY LIMIT



# BASIC TESTING – PACKET DENSITY LIMIT

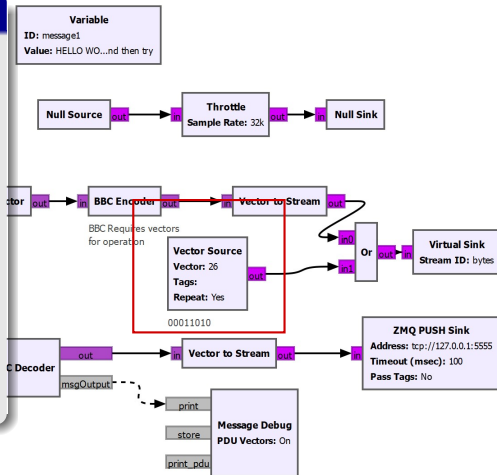
## Details

- Replaced message2 with vector source
- Idea is to test  $\mu = 50\%$  mark density
- $26_{10} = 00011010_2$  (hovering around limit)

PASS

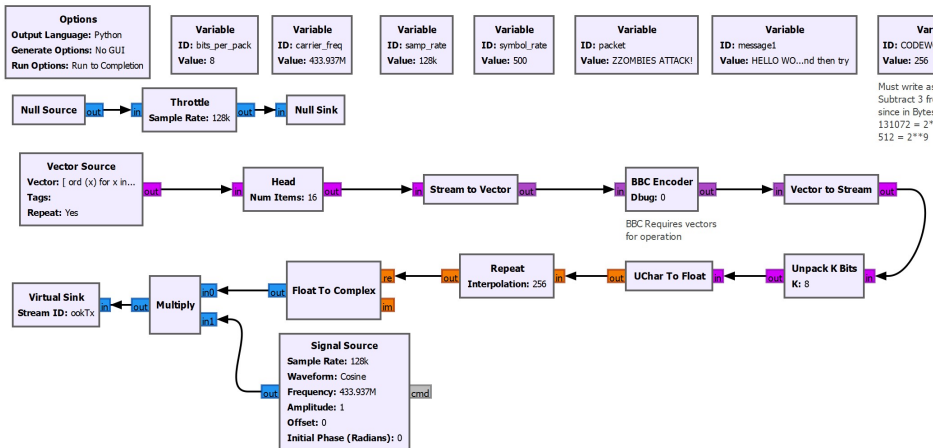
- $27_{10} = 00011011_2$  (over limit)

FAIL



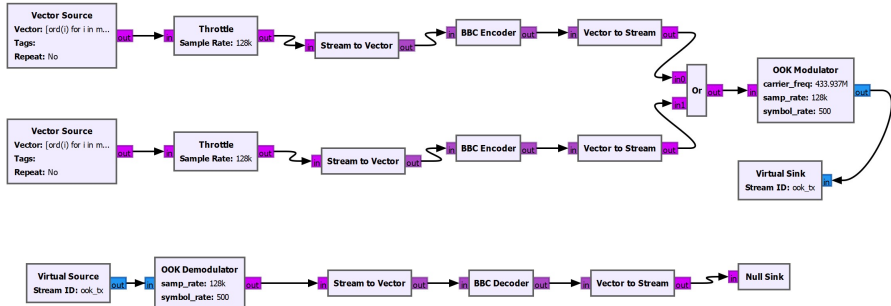
# OOK IMPLEMENTATION

- Chose OOK for simplicity
- Complements jam-resistance of BBC



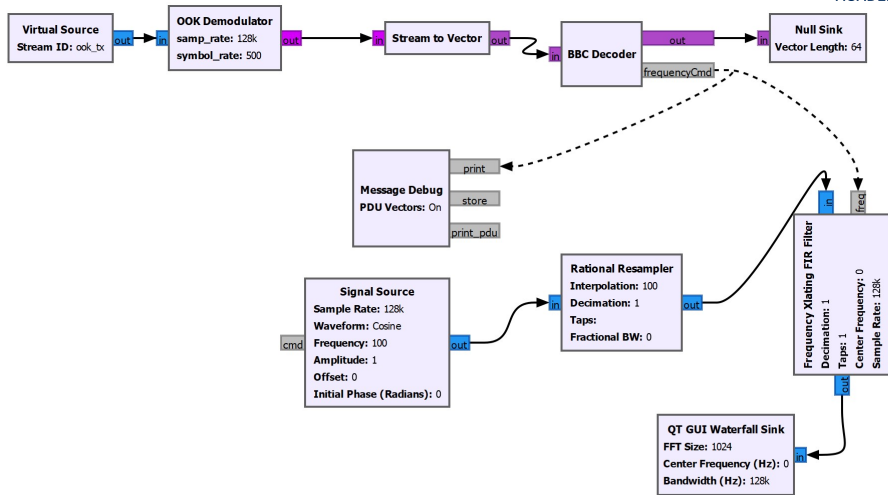
# OOK IMPLEMENTATION

- OOK Hier blocks
- Sends codeword as bytes over message port as PMT.intern()
- BBC Decoder modified to receive bytes/re-assemble codewords
- Outputs messages





# USE CASE - FHSS



# LESSONS LEARNED



- `out_sig` can be a challenge
- Dynamic sizing of in/out vectors
- Sometimes, you just need a PMT
- Block development documentation difficult to follow (but is getting better!)
- Future goal is to output bytes to `out_sig`

# SUMMARY & WHAT'S NEXT?



- Fix bytes output
- Package for CGAN
- Performance benchmark testing
- Hardware testing/implementation
- Active statistical thresholding
- Multimark BBC
- Codeword detection using ring buffer